

# Structural completeness and lattice of extensions in many-valued logics with rational constants

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## 1 Introduction

The logics **RL**, **RP**, and **RG** are obtained by expanding Łukasiewicz logic **L**, product logic **P**, and Gödel logic **G** with rational constants  $\{c_q : q \in [0, 1] \cap \mathbb{Q}\}$  and adding the bookkeeping axioms: For every  $p, q \in [0, 1] \cap \mathbb{Q}$ ,

$$c_p \cdot c_q \leftrightarrow c_{p*q} \quad (c_p \rightarrow c_q) \leftrightarrow c_{p \Rightarrow q} \quad c_0 \leftrightarrow \perp \quad c_1$$

where  $*$  is the Łukasiewicz, Product, and Gödel (minimum) t-norm and  $\Rightarrow$  is the Łukasiewicz, Product, and Gödel standard residuated implication in each case.

The history of these logics goes back to the pioneering works of Goguen [13] and Pavelka [20, 21, 22]. Expanding the language with constants can be viewed as taking advantage of the rich algebraic setting to gain more expressivity; see, e.g., [2, 4, 9, 10, 15, 23, 24]. In this talk, we study the lattices of extensions and structural completeness of these three expansions, obtaining results that stand in contrast to the known situation in **L**, **P**, and **G**.

A rule is an expression of the form  $\Gamma \triangleright \varphi$ , where  $\Gamma \cup \{\varphi\} \subseteq Fm$  is a finite set. A rule  $\Gamma \triangleright \varphi$  is said to be *derivable* in a logic  $\vdash$  when  $\Gamma \vdash \varphi$ . It is *admissible* in  $\vdash$  when for every substitution  $\sigma$  on  $Fm$ ,

$$\text{if } \emptyset \vdash \sigma(\gamma) \text{ for all } \gamma \in \Gamma, \text{ then } \emptyset \vdash \sigma(\varphi).$$

In other words, a rule is admissible in  $\vdash$  when its addition to  $\vdash$  does not produce any new theorem. Clearly, every rule that is derivable in  $\vdash$  is also admissible in  $\vdash$ . If the converse holds,  $\vdash$  is said to be *structurally complete* (SC). Logics whose extensions are all structurally complete have been called *hereditarily structurally complete* (HSC).

During the last two decades, research in structural completeness has turned also to the family of fuzzy logics. While **G** and **P** are hereditarily structurally complete [8, 5], **L** is structurally incomplete [7] and a base for its admissible rules was exhibited by Jeřábek [17], see also [16, 18]. Admissibility in extensions of **L** was investigated in [11, 12].

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Łukasiewicz logic  $\mathbf{L}$ , product logic  $\mathbf{P}$ , and Gödel logic  $\mathbf{G}$  can be obtained as axiomatic extensions of Hájek's basic logic  $\mathbf{BL}$  [14], even if they were defined independently prior to the definition of  $\mathbf{BL}$ . All logics in the  $\mathbf{BL}$  family are algebraizable in the sense of Blok and Pigozzi [3]: the equivalent algebraic semantics of the three logics are the varieties of MV-algebras, product algebras and Gödel algebras, respectively. In fact, from algebraizability we obtain a dual lattice isomorphism from the lattice of finitary extensions of each logic  $L_E(\vdash)$  into the lattice of quasivarieties of each equivalent variety semantics  $L_Q(\mathbb{V})$ . Moreover, if we restrict this isomorphism to the lattice of axiomatic extensions  $L_{AE}(\vdash)$  we get an isomorphism  $L_{AE}(\vdash) \cong L_V(\mathbb{V})$ , where  $L_V(\mathbb{V})$  denotes the lattice of all subvarieties of  $\mathbb{V}$ .

Komori in [19] characterizes  $L_{AE}(\mathbf{L})$  which forms an infinite non totally ordered denumerable pseudo Boolean algebra. The lattice of all extensions  $L_E(\mathbf{L})$  is as complicated as it can be, since the class of all MV-algebras  $\mathbf{MV}$  is Q-universal [1]. That is, for every quasivariety  $\mathbb{K}$  of finite type  $L_Q(\mathbb{K})$  is a homomorphic image of a sublattice of  $L_Q(\mathbf{MV})$ .

The lattice of all axiomatic extensions of  $\mathbf{P}$  is just the three element chain where the only consistent proper axiomatic extension of  $\mathbf{P}$  is classical logic. Since  $\mathbf{P}$  is hereditary structurally complete  $L_E(\mathbf{P}) = L_{AE}(\mathbf{P})$  [6].

Finally, every axiomatic consistent extension of  $\mathbf{G}$  is a finite valued Gödel logic and  $L_{AE}(\mathbf{G}) \cong \omega + 1$ . Since  $\mathbf{G}$  is hereditary structurally complete  $L_E(\mathbf{G}) \cong \omega + 1$  (see [8]).

Structural completeness and the structure of the lattice of axiomatic extensions and the lattice of extensions need not to be preserved when expanding with rational constants, while algebraizability is preserved:

- (i)  $\mathbf{RL}$  is an algebraizable conservative expansion of  $\mathbf{L}$  and the variety of all rational MV-algebras  $\mathbb{RMV}$  is its equivalent variety semantics.
- (ii)  $\mathbf{RP}$  is an algebraizable conservative expansion of  $\mathbf{P}$  and the variety of all rational product algebras  $\mathbb{RP}$  is its equivalent variety semantics.
- (iii)  $\mathbf{RG}$  is an algebraizable conservative expansion of  $\mathbf{G}$  and the variety of all rational Gödel algebras  $\mathbb{RG}$  is its equivalent variety semantics.

We recall that a *rational MV-algebra*, *rational product algebra* and *rational Gödel algebra* is an algebra  $\mathbf{A}$  in the language  $\mathcal{L} = \{\wedge, \vee, \cdot, \rightarrow, \perp, \top\} \cup \{c_q : q \in [0, 1] \cap \mathbb{Q}\}$  such that the  $\{\wedge, \vee, \cdot, \rightarrow, \perp, \top\}$ -reduct is an MV-algebra, Product algebra and Gödel algebra respectively and it satisfies the following bookkeeping equations: For every  $p, q \in [0, 1] \cap \mathbb{Q}$ ,

$$c_p \cdot c_q \approx c_{p \cdot q} \quad (c_p \rightarrow c_q) \approx c_{p \Rightarrow q} \quad c_0 \approx \perp \quad c_1 \approx \top$$

## 2 Main results

### 2.1 Rational Łukasiewicz logic

For the case of Łukasiewicz adding rational constants trivializes the lattice of extensions:

**Theorem 2.1.**  *$\mathbf{RL}$  has no proper consistent extensions, hence  $\mathbf{RL}$  is hereditary structurally complete.*

$$L_E(\mathbf{RL}) = L_{AE}(\mathbf{RL}) \cong 2$$

## 2.2 Rational Product logic

In the case of product logic adding rational constants does not have a significant change in the lattice of axiomatic extensions

**Theorem 2.2.** *RP has two proper consistent axiomatic extensions: namely PL and CL.*

- PL is axiomatized by  $c_q$  for each (some)  $q \in (0, 1] \cap \mathbb{Q}$
- CL is axiomatized by  $c_q$  for each (some)  $q \in (0, 1] \cap \mathbb{Q}$  plus  $\varphi \vee (\varphi \rightarrow \perp)$

**Corollary 2.3.**  $L_{AE}(\mathbf{RP})$  is a four element chain.

Notice that PL is equivalent to the original **P** and CL is equivalent to classical logic, hence when studying admissible rules we will only need to study admissible rules for **RP**.

**Theorem 2.4.** *Every proper extension of RP is structurally complete, but RP is not structurally complete. A base for the admissible rules of RP is given by the set of rules of the form*

$$c_q \vee z \triangleright z \quad (c_p \leftrightarrow x^n) \vee z \triangleright z,$$

for each (equiv. some)  $q \in (0, 1) \cap \mathbb{Q}$  and each  $p \in [0, 1] \cap \mathbb{Q}$ ,  $n \in \omega$  such that  $\sqrt[n]{p}$  is irrational.

Finally, the biggest contrast expanding with rational constants is in the lattice of extensions. We can not obtain a nice description of the lattice  $L_E(\mathbf{RP})$  because of the following result:

**Theorem 2.5.** *The variety RP is Q-universal.*

## 2.3 Rational Gödel Logic

Expanding Gödel logic with rational constants have a significant effect in the lattice of axiomatic extensions. In fact next result shows that we go from a numerable chain to an uncountable chain.

**Theorem 2.6.** *Every consistent axiomatic extension of RG is of the form*

$\mathbf{RG}_r := \mathbf{RG} + \{c_q : q \in [r, 1] \cap \mathbb{Q}\}$  for some  $r \in (0, 1]$ ,

$\mathbf{RG}_p^\omega := \mathbf{RG} + \{c_q : q \in (p, 1] \cap \mathbb{Q}\}$  for some rational  $p \in [0, 1)$  or

$\mathbf{RG}_p^n := \mathbf{RG}_p^\omega + \bigvee_{0 \leq i < j \leq n+2} (c_p \vee x_i) \leftrightarrow (c_p \vee x_j)$  for some rational  $p \in [0, 1)$  and  $n \in \omega$ .

Moreover,  $L_{AE}(\mathbf{RG})$  is an uncountable chain dually isomorphic to the poset obtained adding a new bottom element to the Dedekind–MacNeille completion of the lexicographic order of  $([0, 1) \cap \mathbb{Q}) \times (\omega + 1)$ .

Observe that  $\mathbf{RG}_1 = \mathbf{RG}$  and that  $\mathbf{RG}_0^\omega$  is equivalent to **G** and  $\mathbf{RG}_0^n$  is equivalent to the  $(n + 2)$ -valued Gödel logic. Next result shows that none of the other extensions of **RG** is structurally complete.

**Theorem 2.7.** *The only consistent axiomatic extensions of RG structurally complete are  $\mathbf{RG}_0^\omega$  and  $\mathbf{RG}_0^n$  for each  $n \in \omega$ . Moreover, for all  $r \in (0, 1]$ ,  $p \in [0, 1) \cap \mathbb{Q}$ , and  $\gamma \in \omega + 1$ :*

- A base for the admissible rules of  $\mathbf{RG}_r$  is given by the rules of the form  $c_q \vee z \triangleright z$ , for all  $q \in [0, r) \cap \mathbb{Q}$ ;
- A base for the admissible rules of  $\mathbf{RG}_p^\gamma$  is given by the rule  $c_p \vee z \triangleright z$ .

If we denote by  $\overline{\mathbf{RG}}_r$  the structural completion of  $\mathbf{RG}_r$ , then  $\{\overline{\mathbf{RG}}_r : r \in (0, 1]\}$  is an uncountable antichain in  $L_E(\mathbf{RG})$ . Consequently,  $L_E(\mathbf{RG})$  seems not easy to describe since it contains an uncountable antichain and, by Theorem 2.6, it contains an uncountable chain. The question whether **RG** is Q-universal remains open.

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