## Belief functions over Belnap-Dunn logic

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We have to deal with inconsistent information al the time (media, databases, scientific information, ...)

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(Dunn: "... too much of a good thing")

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Inconsistencies: shortcomings of the available information
(Dunn: "... too much of a good thing")

- issue that has to be lived with until they can be resolved
- many attempts in the literature to devlop paraconsistent logics for these scenarios
- we build our work on the Belnap-Dunn four valued logic
- probability - the most prominent representation of uncertainty
- defined on a boolean algebra
- each value precisely known (subjective probability)
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## Alternatives:

- avoiding precise values (interval probabilities, ...)
- probability measures on weaker structures
- weaker uncertainty measures (inner outer measures, belief functions,... )
(1) Representing incomplete/contradictory probabilistic information
- Belnap-Dunn Logic
- Non-standard probabilities
(2) Dempster-Shafer theory
- Mass functions, belief functions and plausibility functions
- Representation of evidence
(3) Dempster-Shafer theory and BD logic
- Belief functions
- Plausibility?
- Combination of evidence
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## Belnap-Dunn square 4 [Belnap 19]

Belnap-Dunn square $(4, \wedge, \vee, \neg)$ is a de Morgan algebra.

- $(4, \wedge, \vee)$ is a lattice
- each element represents the available positive and/or negative information
- $n$ : no information
- $f$ : false (is bottom)
- $t$ : true (is top)
- b: contradictory information
- $\neg$ is an involutive de Morgan negation.


Belnap-Dunn square 4


Independence of positive and negative information.

Language. $L_{\mathrm{BD}} \ni \varphi:=p \in \operatorname{Prop}|\varphi \wedge \varphi| \varphi \vee \varphi \mid \neg \varphi$

BD Models. $M=\left\langle W, v^{+}, v^{-}:\right.$Prop $\left.\rightarrow \mathcal{P}(W)\right\rangle$
$v^{+}(p)$ : states containing information supporting $p$
$v^{-}(p)$ : states containing information refuting $p$

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Semantics. Two satisfaction relations $\vDash^{+}, \mathfrak{F}^{-}$

$$
\begin{aligned}
& w \vDash^{+} p \text { iff } w \in v^{+}(p) \\
& w \mathfrak{F}^{+} \neg \phi \text { iff } w \vDash^{-} \phi \\
& w \mathfrak{F}^{+} \phi \wedge \phi^{\prime} \text { iff } w ₹^{+} \phi \text { and } w F^{+} \phi^{\prime} \\
& w F^{-} \phi \wedge \phi^{\prime} \text { iff } w \vDash^{-} \phi \text { or } w \vDash^{-} \phi^{\prime} \\
& w \mathfrak{F}^{+} \phi \vee \phi^{\prime} \text { iff } w \mathfrak{F}^{+} \phi \text { or } w \mathfrak{F}^{+} \phi^{\prime} \\
& w \mathfrak{F}^{-} \phi \vee \phi^{\prime} \text { iff } w \mathfrak{F}^{-} \phi \text { and } w \mathfrak{F}^{-} \phi^{\prime}
\end{aligned}
$$

Language. $L_{\mathrm{BD}} \ni \varphi:=p \in \operatorname{Prop}|\varphi \wedge \varphi| \varphi \vee \varphi \mid \neg \varphi$

BD Models. $M=\left\langle W, v^{+}, v^{-}:\right.$Prop $\left.\rightarrow \mathcal{P}(W)\right\rangle$
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& w \vDash^{-} p \text { iff } w \in v^{-}(p) \\
& w \mathfrak{F}^{-} \neg \phi \text { iff } w \vDash^{+} \phi \\
& w \mathfrak{F}^{+} \phi \wedge \phi^{\prime} \text { iff } w \mathfrak{F}^{+} \phi \text { and } w \mathfrak{F}^{+} \phi^{\prime} \quad w \vDash^{-} \phi \wedge \phi^{\prime} \text { iff } w \mathfrak{F}^{-} \phi \text { or } w \mathfrak{F}^{-} \phi^{\prime} \\
& w \mathfrak{F}^{+} \phi \vee \phi^{\prime} \text { iff } w \vDash^{+} \phi \text { or } w \vDash^{+} \phi^{\prime} \quad w \mathfrak{F}^{-} \phi \vee \phi^{\prime} \text { iff } w \mathfrak{F}^{-} \phi \text { and } w \mathfrak{F}^{-} \phi^{\prime}
\end{aligned}
$$

Can we introduce probability in this framework?

## Non-standard probabilities

Frame semantics (Klein, Majer, Rafiee Rad 2021)

- independence of positive and negative probabilistic information
- BD model extended with a (classical) probability measure.


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Frame semantics (Klein, Majer, Rafiee Rad 2021)

- independence of positive and negative probabilistic information
- BD model extended with a (classical) probability measure.

A probabilistic BD model is a tuple $M=\left\langle W, v^{+}, v^{-}, m\right\rangle$, s.t. $\left\langle W, v^{+}, v^{-}\right\rangle$is a BD model and $\mu: \mathcal{P}(W) \rightarrow[0,1]$ is a probability measure on $W$

## Positive probability

$$
p^{+}(\varphi):=\mu\left(\|\varphi\|^{+}\right)
$$

Negative probability of $\varphi: \mathrm{p}^{-}(\varphi):=\mu\left(\|\varphi\|^{-}\right)$
Remark $p^{-}(\varphi)=p^{+}(\neg \varphi)$, otherwise $p^{+}(\varphi)$ and $p^{-}(\varphi)$ independent,

## Theorem [Klein et al]

The following axiomatization of non-standard probabilities is complete with respect to the class of probabilistic frames.
(A1) normalization
$0 \leq p^{+}(\varphi) \leq 1$
(A2) monotonicity
(A3) import-export
if $\varphi \vdash_{B D} \psi$ then $p^{+}(\varphi) \leq p^{+}(\psi)$
$p^{+}(\varphi \wedge \psi)+p^{+}(\varphi \vee \psi)=p^{+}(\varphi)+p^{+}(\psi)$.

## Non-standard probabilities: axioms

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## Remarks

- $p^{-}(\varphi)=p^{+}(\neg \varphi)$
- weaker than classical Kolmogorovian axioms.
(A3 instead of additivity).
- In general $p^{+}(\neg \varphi) \neq 1-p^{+}(\varphi)$
- one can have $0<p^{+}(\varphi \wedge \neg \varphi), p^{+}(\varphi \vee \neg \varphi)<1$


## Non-standard probabilities: intuitive representation

Continuous extension of Belnap-Dunn square, which we can see as the product bilattice $\mathbf{L}_{[0,1]} \odot \mathbf{L}_{[0,1]}$ with $\mathbf{L}_{[0,1]}=([0,1]$, min, max $)$.

- $\left(p^{+}(\varphi), p^{-}(\varphi)\right)$ : positive and negative probabilistic support of $\varphi$.
- $(0,0)$ : no information concerning $\varphi$ is available
- ( 1,1 ): maximally conflicting information
- vertical dashed line: "classical" case
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## Dempster Shafer belief functions

- $\operatorname{bel}(A)$ represents total evidence supporting $A$
- defined on a boolean algebra of events
- weaker than probability
- complex formulas are not determined by the simpler ones $\operatorname{bel}(A \vee B) \geq \operatorname{bel}(A)+\operatorname{bel}(B)$ for $A, B$ disjoint
- provides a lower bound for 'true' probability


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## Example

Scenario. A patient has disease $a, b$ or $c$.
A doctor says "given the evidence she has disease $a$ or $b$ with a high certainty (e.g. 0.7)." However the doctor gives only a low certainty to each of $a, b$ (e.g. 0.1).

Evidence via mass functions

## Mass function

- evidence assigned exactly to a particular event

Definition
A function $\mathrm{m}: \mathcal{P}(S) \rightarrow[0,1]$ is a mass function if

- m(Ø)
- $\sum_{A \in \mathcal{P}(S)} \mathrm{m}(A)=1$.


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Belief function via mass function:

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\operatorname{bel}(A)=\sum_{B \subseteq A} \mathrm{~m}(B)
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$$

Probability via mass function:

$$
p(A)=\sum_{B \subseteq A} m(B)=\sum_{s \in A} m(\{s\})
$$

All information is encoded in singletons.

## Explicit representation of belief functions

## Definition

$f: \mathcal{P}(S) \rightarrow[0,1]$ is a belief function if $f(\emptyset)=0, f(S)=1$ and

$$
\begin{equation*}
f\left(\bigvee_{1 \leq i \leq k} A_{i}\right) \geq \sum_{\substack{J \subseteq\{1, \ldots, k\} \\ j \neq \varnothing}}(-1)^{|J|+1} \cdot f\left(\bigwedge_{j \in J} A_{j}\right) \tag{1}
\end{equation*}
$$

holds for every $k \geq 1$, and for every $A_{1}, \ldots, A_{k} \in \mathcal{P}(S)$.

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## Theorem

For every belief function bel there is a mass function $m_{\text {bel }}: \mathcal{P}(S) \rightarrow[0,1]$ such that, for every $A \in \mathcal{P}(S)$,

$$
\operatorname{bel}(A)=\sum_{B \leq A} \mathrm{~m}_{\mathrm{bel}}(B)
$$

## Plausibility functions

- dual to belief
- $\mathrm{pl}(A)$ represents evidence which is compatibe with $A$
- gives an upper bound for 'true' probability

Plausibility from belief

$$
\operatorname{pl}(A)=1-\operatorname{bel}(\neg A)
$$

Plausibility via mass function:

$$
\operatorname{bel}(A)=\sum_{B \cap A \neq \emptyset} m(B)
$$

Explicit definition

$$
\begin{equation*}
f\left(\bigwedge_{1 \leq i \leq k} A_{i}\right) \leq \sum_{\substack{J \subseteq\{1, \ldots, k\} \\ j \neq \varnothing}}(-1)^{|J|+1} \cdot f\left(\bigvee_{j \in J} A_{j}\right) \tag{2}
\end{equation*}
$$

## Example

Scenario. A patient has disease $a, b$ or $c$.
A doctor says "the patient has disease a or b with certainty 0.7."
The doctor gives no information about disease $c$.

- $\mathrm{m}: \mathcal{P}(S) \rightarrow[0,1]$ is computed based on the evidence
- $\operatorname{bel}(A)=\sum_{B \leq A} m(B)$ : the evidence supporting a
- $\mathrm{pl}(A)=1-\operatorname{bel}(\neg A)=\sum_{B \cap A \neq \emptyset} \mathrm{m}(B)$ : the evidence not contradicting $A$
- $\operatorname{bel}(A) \leq \operatorname{pl}(A)$.


## Representation of evidence. An example

Scenario. A patient has disease $a, b$ or $c$.
A doctor says "the patient has disease a or b with certainty 0.7." It is assumed it is impossible for the patient to have two of them.

## Representation

- $S=\{a, b, c\}$ and m, bel, $\mathrm{pl}: \mathcal{P}(S) \rightarrow[0,1]$
- $\mathrm{m}(\{a, b\})=0.7$ and $\mathrm{m}(S)=0.3$.



## An example

Scenario. A patient has disease $a, b$ or $c$.
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## Representation

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- $\mathrm{m}(\{a, b\})=0.7$ and $\mathrm{m}(S)=0.3$.

We get:

$$
\begin{array}{ll}
\operatorname{bel}(\{a\})=\operatorname{bel}(\{b\})=\operatorname{bel}(\{c\})=0 & \\
\operatorname{bel}(\{a, b\})=\sum X \subseteq\{a, b\} \\
\operatorname{mi}(X)=0.7 & \operatorname{pl}(\{a, b\})=1-\operatorname{bel}(\{c\})=1 \\
\operatorname{pl}(\{a\})=\operatorname{pl}(\{b\})=1 & \operatorname{pl}(\{c\})=1-\operatorname{bel}(\{a, b\})=0.3
\end{array}
$$

- $m(\{a, b\})$ : the 'probability' that the disease is in the set $\{a, b\}$ without being able to say to which subset it belongs.


## Dempster-Shafer combination rule

Let $m_{1}$ and $m_{2}$ be two mass functions on $\mathcal{P}(S)$. Dempster-Shafer combination rule computes their aggregation $\mathrm{m}_{1 \oplus 2}: \mathcal{P}(\mathrm{S}) \rightarrow[0,1]$ as follows.

$$
\mathrm{m}_{1 \oplus 2}(X) \mapsto\left\{\begin{array}{lr}
0 & \text { if } X=\varnothing \\
\frac{\sum\left\{\mathrm{m}_{1}\left(X_{1}\right) \cdot \mathrm{m}_{2}\left(X_{2}\right) \mid X_{1} \cap X_{2}=X\right\}}{N} & \text { otherwise }
\end{array}\right.
$$

Normalization factor:

$$
\begin{aligned}
N & =\sum\left\{m_{1}\left(X_{1}\right) \cdot m_{2}\left(X_{2}\right) \mid X_{1} \cap X_{2} \neq \varnothing\right\} \\
& =1-\sum\left\{m_{1}\left(X_{1}\right) \cdot m_{2}\left(X_{2}\right) \mid X_{1} \cap X_{2}=\varnothing\right\}
\end{aligned}
$$

## Scenario

A patient has disease $a, b$ or $c$.
Doctor 1:
"the patient has a with certainty 0.9 and $b$ with certainty 0.1 ."
Doctor 2:
"the patient has $c$ with certainty 0.9 and $b$ with certainty 0.1 ."

## Representation

$$
\begin{aligned}
& S=\{a, b, c\} \\
& m_{1}(\{a\})=0.9 \text { and } m_{1}(\{b\})=0.1 . \\
& m_{2}(\{c\})=0.9 \text { and } m_{2}(\{b\})=0.1 .
\end{aligned}
$$

DS combination rule ignores contrtadictory information

$$
\mathrm{m}_{1 \oplus 2}(\{b\})=1, \mathrm{~m}_{1 \oplus 2}(\{a\})=\mathrm{m}_{1 \oplus 2}(\{c\})=0
$$

because $\{a\} \cap\{b\}=\{a\} \cap\{c\}=\emptyset$.
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Extending belief with a (classical) belief function

## Definition belief on BD-models

$M=\left\langle W, v^{+}, v^{-}, B e l\right\rangle$, a BD model plus Bel : $\mathcal{P}(W) \rightarrow[0,1]$.

$$
\operatorname{bel}^{+}(\phi):=\operatorname{Bel}\left(|\phi|^{+}\right) \quad \text { and } \quad \operatorname{bel}^{-}(\phi):=\operatorname{Bel}\left(|\phi|^{-}\right)
$$

bel $^{+}$: belief function on the associated Lindenbaum algebra $\mathcal{L}_{\mathrm{BD}}$. bel ${ }^{-}$: belief function on $\mathcal{L}_{\mathrm{BD}}^{o p}$.

Remark. if $\perp$ and $T$ are not in the language bel $^{+}\left(\right.$resp. $\left.\mathrm{pl}^{+}\right)$are general belief (resp. plausibility) functions.

## Non-standard probabilities

Models: $\left(W, v^{+}, v^{-}, \mathrm{m}: W \rightarrow[0,1]\right)$
$p^{+}(\phi)=\sum_{s \in|\phi|^{+}} \mathrm{m}(s)$ and $p^{-}(\phi)=\sum_{s \in|\phi|^{-}} \mathrm{m}(s)$

## Immediate generalisation for belief.

## Non-standard beliefs

Models: $\left(W, v^{+}, v^{-}, m: \mathcal{P}(W) \rightarrow[0,1]\right)$ bel $^{+}(\phi)=\sum_{X \subseteq|\phi|^{+}} \mathrm{m}(X)$ and $\operatorname{bel}^{-}(\phi)=\sum_{X \subseteq|\phi|^{-}} \mathrm{m}(X)$

- bel $^{+}(\phi)$ : belief that $\phi$ is true
- bel $^{-}(\phi)$ : belief that $\phi$ is false
- bel $^{+}$satisfies the axioms of belief functions

Two-dimensional interpretation
Two dimensional reading allows for various combinations of positive/negative belief/plausibility:

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- classically bel $(\phi)=1-\mathrm{pl}(\neg \phi)$. $\mathrm{pl}(\neg \phi)$ maximum evidence against $\phi$ we can consider in two dimensional reading it corresponds to $\mathrm{pl}^{-}(\phi)$

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\left(\mathrm{bel}^{+}(\phi), \mathrm{pl}^{-}(\phi)\right)
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- consider both belief and plausibility independently:

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\left(\operatorname{bel}^{+}(\phi), \operatorname{bel}^{-}(\phi)\right),\left(\mathrm{pl}^{+}(\phi), \mathrm{pl}^{-}(\phi)\right)
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## Two-dimensional interpretation

Two dimensional reading allows for various combinations of positive/negative belief/plausibility:

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- consider both belief and plausibility independently:

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\left(\mathrm{bel}^{+}(\phi), \operatorname{bel}^{-}(\phi)\right),\left(\mathrm{pl}^{+}(\phi), \mathrm{pl}^{-}(\phi)\right)
$$

- If we require $\operatorname{bel}(X) \leq \mathrm{pl}(X)$, for $X \in \mathcal{P}(W)$, then bel and pl must be defined from different mass functions. a piece of evidence might support belief and plausibility in a different way $\rightarrow$ gives rise to two mass functions (e.g., circumstantial evidence vs. direct evidence)


## Combination of evidence

Let $\mathcal{L}$ be a finite distributive lattice.

## Without $\perp$ and $T$

$$
\begin{aligned}
\mathrm{m}_{1 \oplus 2}: \mathcal{L} & \rightarrow[0,1] \\
x & \mapsto \sum\left\{\mathrm{~m}_{1}\left(x_{1}\right) \cdot \mathrm{m}_{2}\left(x_{2}\right) \mid x_{1} \wedge x_{2}=x\right\} .
\end{aligned}
$$

## With $\perp$ and $T$

$$
\begin{aligned}
\mathrm{m}_{1 \oplus 2}: \mathcal{L} & \rightarrow[0,1] \\
x & \mapsto\left\{\begin{array}{lr}
0 & \text { if } x=\perp \\
\frac{\sum\left\{m_{1}\left(x_{1}\right) \cdot m_{2}\left(x_{2}\right) \mid x_{1} \wedge x_{2}=x\right\}}{\sum\left\{m_{1}\left(x_{1}\right) \cdot m_{2}\left(x_{2}\right) \mid x_{1} \wedge x_{2} \neq \perp\right\}} & \text { otherwise. }
\end{array}\right.
\end{aligned}
$$

## Examples. The two doctors

Scenario. A patient has disease $a, b$ or $c$.
Doctor 1 : a with certainty 0.9 and $b$ with certainty 0.1 .
Doctor 2: $c$ with certainty 0.9 and $b$ with certainty 0.1 .
Representation. $\mathrm{m}_{1}, \mathrm{~m}_{2}: \mathcal{D M}_{3} \rightarrow[0,1]$

$$
\mathrm{m}_{1}(x)=\left\{\begin{array}{ll}
0.9 & \text { if } x=a \\
0.1 & \text { if } x=b \\
0 & \text { otherwise }
\end{array} \quad \mathrm{m}_{2}(x)= \begin{cases}0.9 & \text { if } x=c \\
0.1 & \text { if } x=b \\
0 & \text { otherwise }\end{cases}\right.
$$

Dempster-Shafer combination rule gives

$$
\mathrm{m}_{1 \oplus 2}(x)= \begin{cases}0.81 & \text { if } x=a \wedge c \\ 0.09 & \text { if } x=a \wedge b \text { or } x=b \wedge c \\ 0.01 & \text { if } x=b \\ 0 & \text { otherwise }\end{cases}
$$

$\operatorname{bel}_{1 \oplus 2}(a)=\operatorname{bel}_{1 \oplus 2}(c)=0.9$ and $^{b^{\prime}}{ }_{1 \oplus 2}(b)=0.19$

## Examples. The two doctors

Representation. $\mathrm{m}_{1}, \mathrm{~m}_{2}: \mathcal{D}_{3} \rightarrow[0,1]$

$$
\begin{aligned}
& \mathrm{m}_{1}(x)= \begin{cases}0.9 & \text { if } x=a \wedge \neg b \wedge \neg c \\
0.1 & \text { if } x=\neg a \wedge b \wedge \neg c \\
0 & \text { otherwise } .\end{cases} \\
& \mathrm{m}_{2}(x)= \begin{cases}0.9 & \text { if } x=\neg a \wedge \neg b \wedge c \\
0.1 & \text { if } x=\neg a \wedge b \wedge \neg c \\
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\end{aligned}
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$$
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$$

$\operatorname{bel}_{1 \oplus 2}(a)=\operatorname{bel}_{1 \oplus 2}(c)=0.9$ and bel $_{1 \oplus 2}(b)=0.19$

## Reasoning with inconsistent / incomplete uncertain information

- other uncertinty measures (upper/lower probabilities, ...)
- qualitative probability
- various aggregation methods
- two layered framework
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