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Kripke-definable ordinals

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O. It is known (see e.g. [4]) that in propositional K.(ripke)semantic it holds:

PROP.O.0 For every $n \geq 1$ there is a (negation-free) formula $W_n (= p_0 \vee \bigvee_{i=0}^{n-1} (p_i \rightarrow p_{i+1})) = T_{n-1}$ of [11] with the property: it is valid in every K.model on the poset $m = \overset{0}{\cdot} \rightarrow \overset{1}{\cdot} \rightarrow \dots \rightarrow \overset{m-1}{\cdot}$ with $m \leq n$, and has a K.countermodel on $m+1 = \overset{0}{\cdot} \rightarrow \overset{1}{\cdot} \rightarrow \dots \rightarrow \overset{n}{\cdot}$.

But also (as a consequence of the K.characterization with f.m.p. for Dummett's logic LC):

PROP.O.1 For no ordinal $\xi \geq \omega$ there is a formula W_ξ with the above property (w.r. to ξ).

Our aim is to investigate on this kind of connection between ordinals and formulas in first order K.semantic.

1. We use a first order language L with denumerably many n-place predicate symbols P_0^n, P_1^n, \dots for $n \geq 0$. A K.frame is a triple $F = \langle P, \leq, V \rangle$ where (P, \leq) is a poset and V a domain function on it. K.realizations q for L on F and the forcing relation \models are defined as usual. L^*F is the set of all L-sentences α s.t. $q \models \alpha$ for every q on F; \underline{L}^*C (C a class of K.frames) is $\bigcap_{F \in C} L^*F$. By $\mu, \nu, \xi \dots$ ($m, n, s \dots$) we denote ordinals (finite); by \leq their usual order relation. If $\mu \neq 0$, (μ, \leq) is the poset of all

$\nu < \mu$ ordered by \leq . Further, for $\xi > 0$:

- $[\xi;] = \{ \langle P, \leq, V \rangle \mid \exists \mu < \xi. (P, \leq) \cong (\mu, \leq) \}$.

- $[\xi; \text{Con}] = \{ \langle P, \leq, V \rangle \in [\xi;] \mid V \text{ is constant} \}$.

- $[\prec \xi;] = \bigcup_{0 < \mu < \xi} [\mu;]$, $[\prec \xi; \text{Con}] = \bigcup_{0 < \mu < \xi} [\mu; \text{Con}]$.

Recall the special case of Cantor Normal Form Thm.:

PROP. 1.0 $\forall \xi (0 < \xi < \omega^\omega \rightarrow \exists ! n, m_1 \dots m_n \geq 0, m_0 \geq 1. \xi = \omega^n \cdot m_0 + \omega^{n-1} \cdot m_1 + \dots + \omega^0 \cdot m_n)$. In short: $c(\xi) = n; m_0, m_1 \dots m_n$.

2. Looking at Prop.0.0, we say:

DEF. 2.0 An ordinal $\xi > 0$ is K.definable iff there exists a L-sentence W_ξ s.t. $W_\xi \in \underline{L}^*[\xi; 1] \setminus \underline{L}^*[\xi+1; 1]$.

Our first claim is a positive one (compare it with Prop.0.0):

THEOREM 2.1 If $0 < \xi < \omega^\omega$ then ξ is K.definable.

The suitable W_ξ 's are defined as follows:

(a) For $n \geq 1$ and $k \geq 0$, A_k^n is the L-sentence:

$$\forall x_1 \left[\begin{array}{l} \left(\forall x_2 \dots x_n (P_k^n x_1 \dots x_n \rightarrow \forall x_1 \dots x_n P_k^n x_1 \dots x_n) \rightarrow \forall x_1 \dots x_n P_k^n x_1 \dots x_n \right) \\ \wedge \\ \left(\forall x_2 \dots x_n (P_k^n x_1 \dots x_n \rightarrow \forall x_2 \dots x_n P_k^n x_1 \dots x_n) \rightarrow \forall x_2 \dots x_n P_k^n x_1 \dots x_n \right) \\ \wedge \\ \left(\forall x_2 x_3 \dots x_n (P_k^n x_1 \dots x_n \rightarrow \forall x_2 \dots x_n P_k^n x_1 \dots x_n) \rightarrow \forall x_2 \dots x_n P_k^n x_1 \dots x_n \right) \\ \wedge \\ \vdots \\ \wedge \\ \left(\forall x_2 x_3 \dots x_n (P_k^n x_1 \dots x_n \rightarrow \forall x_n P_k^n x_1 \dots x_n) \rightarrow \forall x_n P_k^n x_1 \dots x_n \right) \end{array} \right] \rightarrow \forall x_1 \dots x_n P_k^n x_1 \dots x_n .$$

Note that A_k^1 ($k \geq 0$) is instance of schema H_2^0 considered in [1]: $\forall x ((\alpha(x) \rightarrow \forall y \alpha) \rightarrow \forall y \alpha) \rightarrow \forall x \alpha$.

(b) For all $n, m_1 \dots m_n \geq 0$ and $m_0 \geq 1$, $W[n; m_0, m_1 \dots m_n]$ is the L-sentence (definition by induction on n and subinduction on m_0):

$$W[0; m_0] : P_0^0 \vee \bigvee_{i=0}^{m_0-1} (P_1^0 \rightarrow P_{i+1}^0)$$

$$W[1; m_0, m_1] : \begin{cases} m_0=1 : \begin{cases} A_0^1 \vee (\forall x P_0^1 x \rightarrow P_0^0) , & \text{if } m_1=0 \\ A_0^1 \vee (\forall x P_0^1 x \rightarrow W[0; m_1]) , & \text{if } m_1 > 0 \end{cases} \\ m_0=k+1, k \geq 1 : A_k^1 \vee (\forall x P_k^1 x \rightarrow W[1; k, m_1]) \end{cases}$$

$$W[n+1; m_0, m_1 \dots m_{n+1}] : \begin{cases} m_0=1 : \begin{cases} A_0^{n+1} \vee (\forall x_1 \dots x_{n+1} P_0^{n+1} x_1 \dots x_{n+1} \rightarrow P_0^0) , & \text{if } m_1 = \dots = m_{n+1} = 0 \\ A_0^{n+1} \vee (\forall x_1 \dots x_{n+1} P_0^{n+1} x_1 \dots x_{n+1} \rightarrow W[n+1-s; m_s \dots m_{n+1}]) , & \text{if } m_1 = m_2 = \dots = m_{s-1} = 0 \text{ and } m_s > 0 \end{cases} \\ m_0=k+1, k \geq 1 : A_k^{n+1} \vee (\forall x_1 \dots x_{n+1} P_k^{n+1} x_1 \dots x_{n+1} \rightarrow W[n+1; k, m_1 \dots m_{n+1}]) \end{cases}$$

(c) Finally, for $0 < \xi < \omega^\omega$, we define (by Prop.1.0): $W_\xi = W[c(\xi)]$ (note that negation is not used!).

Then it is shown (by induction on $c(\xi)$):

LEMMA 2.1.1 $W_\xi \in \underline{L}^*[\xi; 1] \setminus \underline{L}^*[\xi+1; \text{Con}]$ (which is stronger than $W_\xi \in \underline{L}^*[\xi; 1] \setminus \underline{L}^*[\xi+1; 1]$ requested by Def.2.0).

Our second claim is a negative one:

THEOREM 2.2 No $\xi \geq \omega_1$ (= first uncountable) is K.definable.

This easily follows from

PROP. 2.2.1 For every frame $F = \langle P, \leq, V \rangle$ in which (P, \leq) is well ordered there exists a subframe $F' = \langle P', \leq', V' \rangle$ in which (P', \leq') is well ordered, $|P| \leq \aleph_0$, $\bigcup_{i \in P} |V_i| \leq \aleph_0$, and s.t. $\underline{L}^* F' \subseteq \underline{L}^* F$.

which in turn is a corollary to a strong Löwenheim Skolem-type Thm. for K.semantic we proved in [2] (improving Ono's result in [3]).

Unfortunately we lack answers for $\{ \xi \mid \omega^\omega \leq \xi < \omega_1 \}$ in terms of K.definability, although we see some reasons for supporting the following

CONJECTURE 2.3 No $\xi \succ \omega^\omega$ is K.definable.

But the proof (if one!) seems to be not so easy.

REMARK 2.4 Say: $\xi(\succ 2)$ is K^o.definable iff there exists a L-sentence $W_\xi^o \in \underline{L}^* \prec \xi; 1 \sim \underline{L}^* \prec \xi + 1; 1$. It can be shown (modifying the W's of Thm.2.1) that every $\xi \prec \omega^\omega$ is K^o.definable and that no $\xi \succ \omega_1$ is such.

REMARK 2.5 From Thm.2.1 we obtain (for notations see [1]):

(a) There exists a strictly decreasing well ordered chain of negation-free logics (finitely axiomatizable?) between TQ and BQ, with order type $\omega^\omega + 1$.

(b) There exists a strictly decreasing well ordered chain of negation-free logics containing schema D (fin.ax.?) between TQ and $BH_1 D$, with order type $\omega^\omega + 1$.

We can get other results studying intermediate logics obtained from a given basis by adding W_ξ of Thm.2.1 as new axiom schema for some ξ . It turns out, e.g., that $PW_\xi \subset PH_2^o$ for all $\omega \preceq \xi \prec \omega^\omega$.

References

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